Automatic Security Evaluation and (Related-key) Differential Characteristic Search: Application to SIMON, PRESENT, LBlock, DES(L) and Other Bit-oriented Block Ciphers

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I. Motivation

2. Existing methods for differential characteristic search

3. Our method

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Motivation

Differential cryptanalysis [Biham, Shamir, 1991] is one of the most powerful attacks on block ciphers

Finding a good differential characteristic with high probability is the first step in the (related-key) differential attack

Existing methods (I)

Matsui's Algorithm

- Mitsuru Matsui, On correlation between the order of S-boxes and the strength of DES, Eurocrypt 1994.
- Branch and Bound approach
- Original method only applicable in the single-key setting

□Variants of Matsui's Algorithm

- Alex Biryukov, Ivica Nikolic.: Search for related-key differential characteristics in DES-like ciphers. FSE 2011
- Brach and Bound approach
- Applicable in the related-key setting, but only for <u>linear key schedule</u> <u>algorithm</u>

Integer programming based method

Existing methods (II)

Integer programming based method

- applicable both in single-key and related-key settings
- can be used to obtain security bounds (bounds of the minimum number of active S-boxes) with respect to differential attack

 \checkmark can not be used to obtain good characteristic directly

 not applicable to bit-oriented block ciphers such as PRESENT, SIMON, DES(L), etc.

• Nicky Mouha, Qingju Wang, Dawu Gu, Bart Preneel. Differential and linear cryptanalysis using mixed-integer linear Programming. Inscrypt 2011.

• Shengbao Wu, Mingsheng Wang. Security Evaluation against Differential Cryptanalysis for Block Cipher Structures, IACR ePrint 2011/551.

Our method: mixed-integer programming based

Integer programming based method

- applicable both in single-key and related-key settings
- can be used to obtain security bounds (bounds of the minimum number of active S-boxes) with respect to differential attack
- \checkmark can be used to obtain good characteristics directly
- applicable to bit-oriented block ciphers such as PRESENT, SIMON, DES(L), etc.

Mixed-integer programming: An example

Mixed-integer linear programming (MILP), an example

✓ Objective function

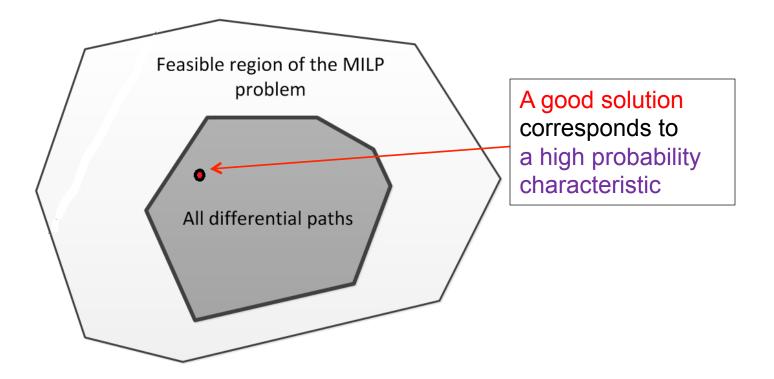
✓ Feasible region: all solutions satisfy the constraints

$$\begin{aligned} \min -x_1 + x_2 - 2x_3 + x_4 - x_5 \\ \text{subject to} \\ x_1 + x_2 &\leq 1 \\ x_1 - 5x_2 + x_3 &\leq 2 \\ 2x_3 + 2x_4 - 4x_5 &\leq 1 \\ x_2 - 2x_4 + x_5 &\leq 0 \\ x &\in \{0, 1\}^5 \end{aligned}$$

Our method: The main idea

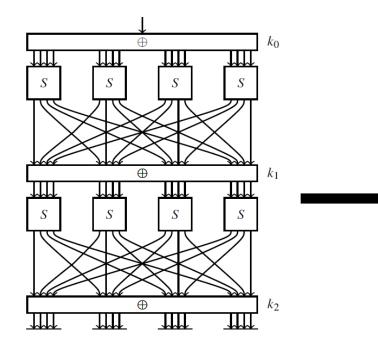
The main idea of our method

- describe the differential behavior of a cipher "at bit-level" by a set of linear inequalities
- try to find a characteristic with minimum number of active S-boxes
- \square Search for high probability characteristic \rightarrow Extract a good solution from the feasible region of an MILP problem

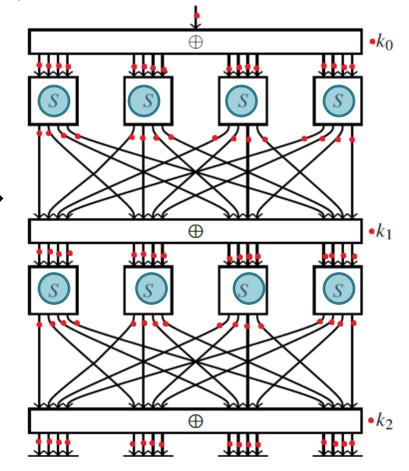


Our method: Modeling technique Variables involved in our MILP model

- for every S-box we introduce a new 0-1 variable (represented by a), such that 1 for active and 0 for otherwise
- for every input and output bit-level difference of every operation we introduce a new 0-1 variable (represeted by a



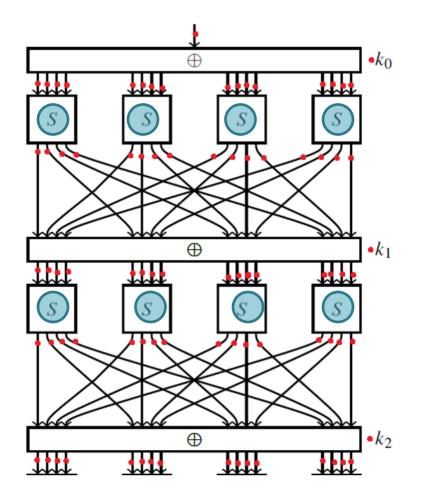
* This toy cipher is taking from Lars R. Knudsen et al's book: *The block cipher The Block Cipher Companion*".



Our method: Modeling technique Objective function \checkmark Minimize the sum of the varaibles (represented by \bigcirc), that is, minimize the number of active S-boxes. Constraints \checkmark Linear inequalities in the variables represented by lacksquare . $\bullet k_0$ \oplus \oplus k_0 S S S S S \oplus k_1 \oplus $\bullet k_1$ S S S S k_2 Ð $\bullet k_2$ \oplus

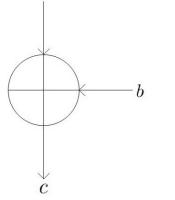
Our method: Constraints generation

How to describe the constraints imposed (by different operations) on the variables denoted by and with linear inequalities ?



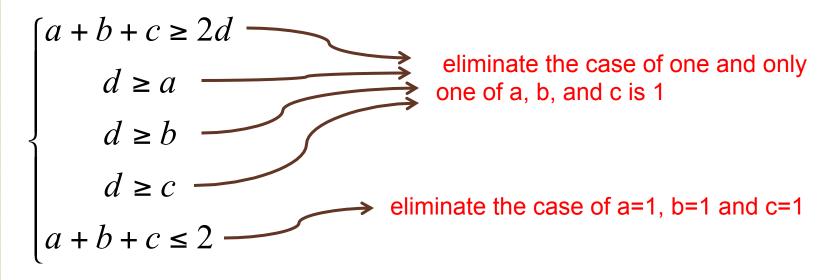
Our method: Constraints generation for XOR

Constraints imposed on the input and output differences by XOR



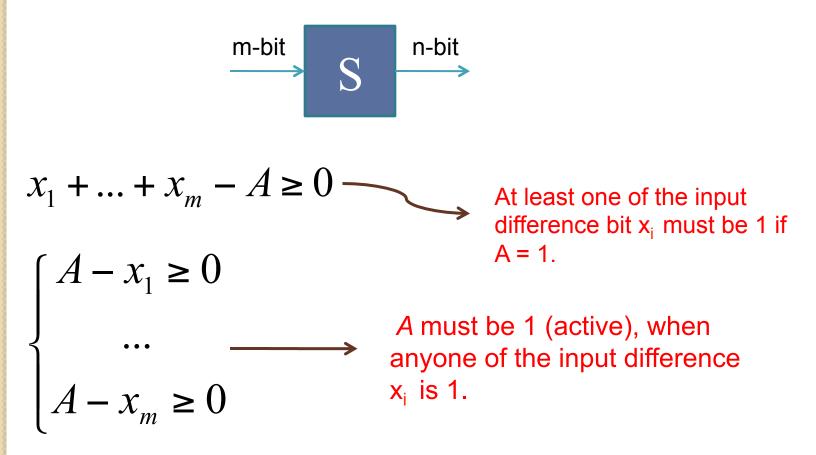
$$a = 0, b = 0 \rightarrow c = 0$$
$$a = 0, b = 1 \rightarrow c = 1$$
$$a = 1, b = 0 \rightarrow c = 1$$
$$a = 1, b = 1 \rightarrow c = 0$$

Constraints (where d is a dummy variable and all variables are 0-1)



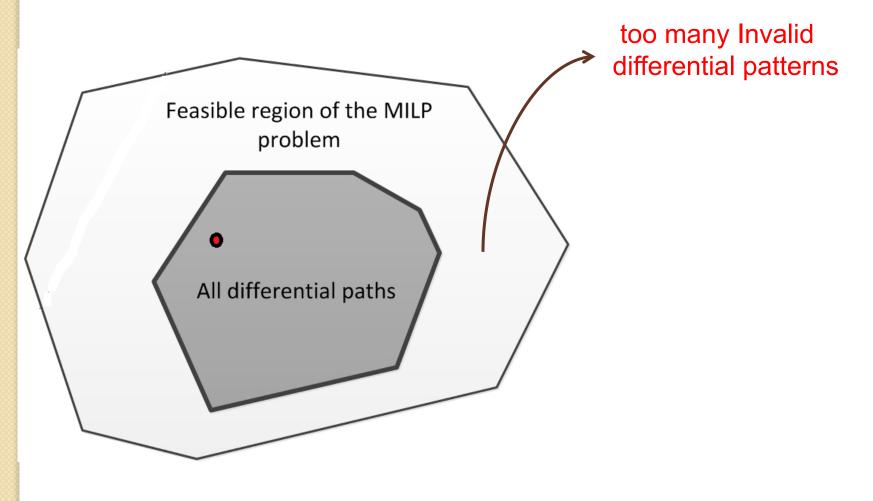
Our method: Constraints generation for S-box

- Constraints imposed on the input and output differences by an m×n S-box (not necessarily invertible)
 - ✓ Let $x_1, x_2, ..., x_m$ be the input difference, and $y_1, y_2, ..., y_n$ be the output difference
 - \checkmark Let A be the variable indicating the activity of the S-box



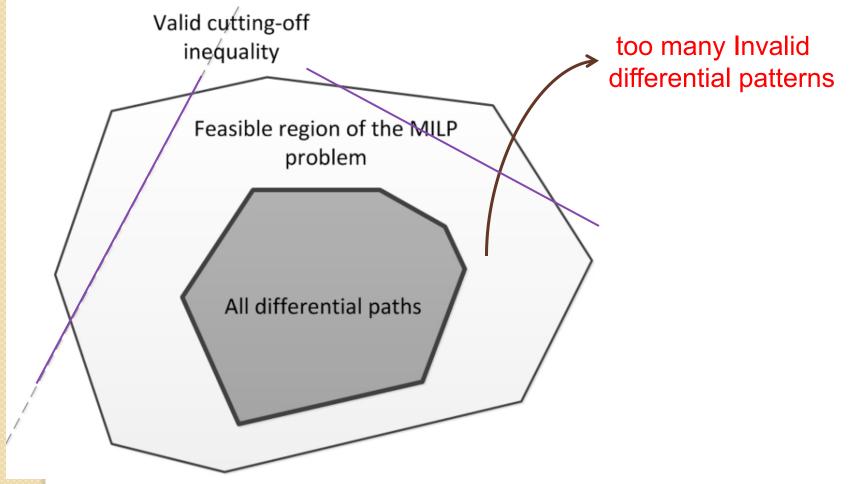
Our method: a more accurate constraints generation

However, this is too coarse to describe an specific S-box, and result in an feasible region contain many invalid differential patterns



Our method: a more accurate constraints generation

Hence, we need the so called valid cutting-off inequalities to remove some impossible differential patterns of an specific S-box.



Our method for constraints generation

Two methods for generating valid cutting-off inequalities for an specific Sbox

I. Logical condition modeling

2. Convex hull computation

Method I

Logical condition modeling

Assume x, y are 0-1 variables, how to describe the logical condition "x must be 1 when y = 1"?

$$x - y \ge 0$$

The differentials of some S-boxes has similar properties.
 For example, the PRESENT S-box.

Fact 1. The S-box of PRESENT-80 has the following properties:

(i) $1001 \rightarrow ***0$: If the input difference of the S-box is 0x9 = 1001, then the least significant bit of the output difference must be 0;

(ii) $0001 \rightarrow ***1$ and $1000 \rightarrow ***1$: If the input difference of the S-box is 0x1 = 0001 or 0x8 = 1000, then the least significant bit of the output difference must be 1;

(iii) ***1 \rightarrow 0001 and ***1 \rightarrow 0100: If the output difference of the S-box is 0x1 = 0001 or 0x4 = 0100, then the least significant bit of the input difference must be 1; and

(iv) *** $0 \rightarrow 0101$: If the output difference of the S-box is 0x5 = 0101, then the least significant bit of the input difference must be 0.

Method I

Logical condition modeling

This conditional differential properties can be described by

$$-x_{0} + x_{1} + x_{2} - x_{3} - y_{3} + 2 \ge 0$$

$$\begin{cases} x_{0} + x_{1} + x_{2} - x_{3} + y_{3} \ge 0 \\ -x_{0} + x_{1} + x_{2} + x_{3} + y_{3} \ge 0 \end{cases}$$

$$\begin{cases} x_{3} + y_{0} + y_{1} + y_{2} - y_{3} \ge 0 \\ x_{3} + y_{0} - y_{1} + y_{2} + y_{3} \ge 0 \end{cases}$$

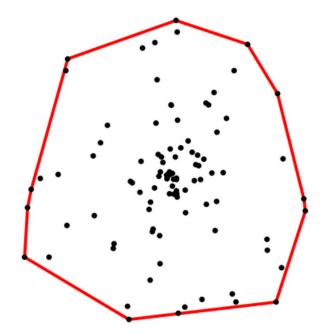
$$-x_{3} + y_{0} - y_{1} + y_{2} - y_{3} + 2 \ge 0$$

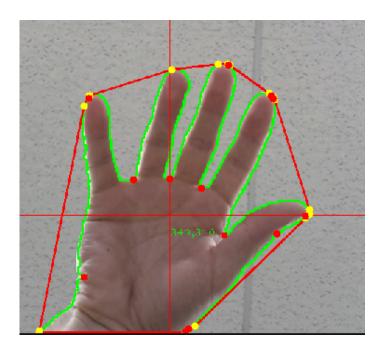
Remove all differential patterns which do not satisfy the differential pattern: $1000 \rightarrow ***0$, $0001 \rightarrow ***1$, $***1 \rightarrow 0001$, $***0 \rightarrow 0101$

Method II

Convex hull computation

 Convex hull of a set of points in Rⁿ: the smallest convex set that contains these points.





Method II

Convex hull computation

 \checkmark A convex hull can be represented by a set of linear inequalities

Treat the set of all possible differential patterns of an S-box as a set of points in Rⁿ. For example, the PRESENT S-box: {(0, 0, 0, 0, 0, 0, 0), (0, 0, 0, 1, 0, 0, 1, 1), (0, 0, 0, 1, 0, 1, 1, 1), (0, 0, 0, 0, 1, 1, 0, 0, 1), (0, 0, 0, 1, 1, 0, 0, 0, 1, 1, 1), (0, 0, 1, 0, 0, 0, 1, 1), (0, 0, 1, 0, 0, 1, 1, 0), (0, 0, 1, 0, 0, 1, 0, 1), (0, 0, 1, 0, 0, 1, 0, 1), (0, 0, 1, 0, 0, 1, 0, 0), (0, 0, 1, 0, 0, 1, 1, 0, 1), (0, 0, 1, 0, 0, 1, 0, 1), (0, 0, 1, 0, 0, 1, 0, 0), (0, 0, 1, 0, 0, 1, 0, 0, 1, 0, 0), (0, 0, 1, 0, 0, 1, 0, 0), (0, 0, 1, 0, 0, 1, 0, 0), (0, 0, 1, 0, 0), (0, 0, 1, 0, 0), (0, 0, 1, 0, 0), (0, 0, 1, 0, 0), (0, 0, 1, 0, 0), (0, 0, 1, 0, 0), (0, 0, 1, 0, 0), (0, 0, 1, 0, 0), (0, 0, 1, 0, 0), (0, 0, 1, 0, 0), (0, 0), (0, 0, 0), (0,

Corresponds to the differential: $0010 \rightarrow 1101$

Then we can compute the linear inequalities representation of the set of differential patterns

Linear inequality description of the PRESENT S-box.

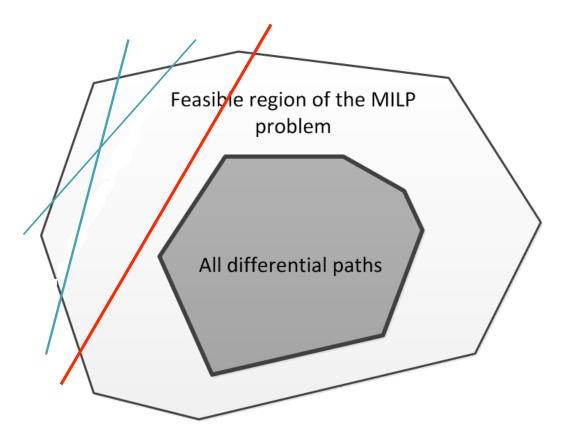
Too many inequalities, which will make the MILP problem too difficult to be solved in practical time

0, 1, 1. 1, 1, -1, 0, 0) (0, 1, 1, 1, -1, 0, 0, 1, 0) (-1, -1, ٥. 1, -1, 0, -1, -1, 4)(1, 0, -1, 2, 1. 2. 2. 01 -1, -1, ٥. 0, -1, 1, -1, 3)(-1, 0, -1, -1, 1, 0, -1, 0, 3)(0, 1, 1. ٥. -1. 1, -1, 1. 1)(1, 1. 1. ٥. ٥. ٥. -1. 1. (1. 1. 2. 0. 1. -1. -1. 0)(-1, 0, 0. 1, 1, 1, 1, 0, 0)(-1, 0, 0. 0)(0, -1. -1. 21 (2. 2. ٥. 1, 1, 1. 1. 1. 1. 1. -1. 0. -1, -2, -2, -2, -1, 6)(0, 1, -1, 2, -1, 3, 2, -1, 0)(1, 0, 1, 1)(1, 2, 0. -1. 0. 0. ٥. ٥. ٥. 0. -1. -1, -1, ٥, 1. 3)(0, 0. -1. ٥. 1. 1. 1. 00 1, -2, -2, -1, -2, -2, 7) (2, 2, -1, 2, -1, 1. -1. 0. 1. 0)(0, 1. -1. 0) (0. 1. 1. 1. 2. 1. -1. 1, -1, 1, 2)(0, 1, 0)(-2, 01 (-2. 1. 1. 3. 2. 0) (1, -1, -1, 1. 1. ٥. 0, -1, 1. -1. 1. 1. 1. ٥. 1. 1. 2. 1. ٥. 1. 1. -1. -1. -1. -1. 1. ٥. -1. 4) (-1, 1, -2, -1, -2, -2, 1. 2, 6) (2, -2, 3, -4, -1, -4, -4, 11)(1, -1, 1. 0) C 0. 1. 2. 2. 2. 0. -1. 3)(0, 1) (1, -1, 2, 2, 2, -1. 1. -1. ٥. -1. -1. 1. 1. -1. 1. ٥. 1. 1. -1. -1. ٥, 2. 0)(-1.٥. -1. 40 -1. -1. ٥. -1. 1, -1, 1. 3) (2, 1, 1. ٥, -2, -1, -2, -1, 4)(1, 0, -1, 1, ٥, 1, 1, -1. 1)(1, 2, 2. ٥. ٥. 1. 01 ٥. ٥. ٥. ٥. 0, -1, ٥. ٥. 1)(1, 1. 1. 1. 1, 1, 1, -2, 0)(1, 0, 1. -1, 1. 1. ٥. 1. 0)(0, -1. 1. -1. -1. 4) 1. 1. -2, 1. 1. ٥. 2. 0) (-2. 1. -1. -1. 1. -1, -1, -1, 5) (1, -1, ٥. 2. 2. 2, 1, -1. 0)(3, 2. 2. ٥. ٥. 1. 1, -1, 1. 0) (0, -1, -1, -1. 0, -1, -1, 4) (2, 2, 2, -1, 1. ٥. -1. 0) (1, -1, -1, ٥. -1. 0. -1. 3) 1. 1. 1. 1. -1. 1. -1. 3) (-2, 0, 0, 2, 1, 0)(1,-1, 1, 2,-2,-2, 1)(1,-1,-1, 1, 2, 0, -2-5)(1, -1, 0, 21 -1. -1. ٥. ٥. 1. 2. 1. 1. 2. ٥. 1. 1. -1. -1. 2. -1. -1. 0)(1, 1. -1. 1, -1, 2, 1, -1, 1. 0)(1, -1, -1, ٥. 2) з. з. з. 1. ٥. 1. -1. -1, 1. 1. 1. 1. 1, 1, -2, 1. 0) (2. 1, 2, -1, 1, 0, -1, 2, 0) (0, -1, -1, -1, -1, 0. 1. -1. 4) (-2. 1. -2. -1. 1. -2. ٥. 6) -1, ٥. 2. -1, -2, -2, -2, 7) (-1, 0, -1, -1, 1, -1, 0, -1, 4) (3, -3, -2, 1, -1, -1, 3, -3, 7)(1, 1, -1, ٥. 0) -1. 1. 1. 1. ٥. 1. -1. -1. -1, -1, ٥. ٥. 3) (1, -2, -2, 0, 1, -2, 1, -1, 5) (-1, -1, -1, -1, з. 4, 3, 4. 0)(2, 2. 2. ٥. 0) ٥. 2 2. -1, 2. 1, -1, 2. 0) (1, -1, -2, -1, -2, 0, 2, -2, 6) (1, -1, 0, -1, 1. 0, -1, -1. 3)(1. ٥. -1. -1. ٥. -1. 1. -1. 3) C 1. 1. ٥. 0, -1, ٥. 0)(1, 0, 1. 1, 1, 1, 0, -1, 0)(0, 0, 0. -1. 0. 0. 0. ٥. 1) (1, -2, -2. -1. 2. -3, 1. -2. 7) 1. ٥. -1. 1. 1. 1. 0) (-1. 1. 1. -2, -1, -1, -1, -2, 6) (1, 2, -1, -2, -2. -1. 1. 1. 4) (2, 2. -1. 1. -1. 2. ٥. 1. ο. 1. -1. 1. ٥. -1. 1)(0. 0. 0. -1. 1. 1. 1. 1. 0)(1. 1. -1. 0. 0. 1. 1. 1. 0) (1, -1, 2. -2. 1. -1. 40 (1. 1. 1. -1, ٥. 1) (-3, 2, -2, -1, 1, -2, -2, -1, 8)(0, ٥, -1, 4)(0, 0. 1. -1. ٥. 1. 1. -1. -1. -1. -1-0. 1. 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(1, -1, 1, 0) (2, ٥. 0, 1, 1. 1. 0) (1, -1, 1, 0, 1, 1, -1, ٥. 1)(-4. 1. 1. 2. з, 2. з. 1. 1. 1. -3. 1. 2. 1 0) C 0. 0, -1, 1, 1, 1. 1. ٥. 0)(-1, 0, -1, -1. 1, 1, 0, 1. 2)(0, 0, -1. 1. 1. ٥. 1. 1. 0)(0. 1. -1. 1. -1. ٥. -1. 3) -1. 0, -1, 1. -1. ٥. -1. 3) (-1, 1, 0, -1. 0, -1, -1, -1, 4) (0, -1, ٥. ٥. ٥, ٥. ٥, ٥. 1)(0, -1. -1. -1. -1. 4) 1. -1. 2. 1. -2. -2. -2. 5) (1, 1, -1, -1, -1, 0, 1, 0. 2) (-1, 1, ٥. -1, 1. -1. 0. -1. 3)(0, -1, ٥. ٥, 1. 1. 1. 1. 0) 1. -1. -1. 0. -1. 3) (0, 1, -1, 4) (2, -3, -1, 2, 8)(0, 0. 1. -1. ٥. -1, -1, -1, -1, -1. 1. -3, -3. ٥. ٥. ٥. ٥. ٥. ٥. -1. 1) (1. 2, -1, -2, -2, -2, -1, ٥. 6) (-1, 1, 1, 2, 0, -1, 1, 2. 0)(1, 1, -2, 1. 1. 1. 1. 1. 0)(1, -1, 1. 1. 1. 1. ٥. 0) -1, -1, -1, 1, -2, 1, -2, 5) (1, 3, -2, -2, 3, 4. 4. 0) (-1, -1, -1, -1, 4. 3)(-2, ٥. 1. 0. 1. 1. 2. -2. -1. -1. 1. -2. 6) C 0 -2, ٥, 1, 2, 1. 2, 0) (1, -1, -2, -2 1, -3, 2, -2, 7) (0, -2, -2, 4, 1. 0)(0, ٥. 2) (1. 1. з. 1. ٥. -1. 1. -1. 1. -1. ٥. 1, 0, -1, 1, 1)(0, 2, 2, 1, -1, 1, -1, 0)(1, 2, -1, -2, -2, -2, 1. 1. 0. 0. 1. 0)(-1, 0, 1. 1. 0. 5) (-1. 1. 2. 1. 2, (-1. 1. 1. -1. 1, -1, 1, -2, 3) (1, 1, 0. 1, -1, 1, 1, ٥. 0)(2, 0. -1. -1. 1. 2. 2. 0)(2, 1. -2. -2. -1. -1. 1. -1. 5) C 4. 1. з. -2. з, 1, -2, 4. 0)(1, 0, 0, -1, -1, -1, -1, 1. 3)(0, 0, -2. 1. 2, 1. 2, 1. 0)(0, ٥. ٥. ٥. -1. ٥. ٥. ٥. 1) 2. 2. 1. 1. -3, 1. 1. 0) (1. 1. ٥. -2, 1. 2, 1. 2. 0)(0,-1, 1. 1. 1. 1. ٥, -1, 1)(-1, 1. ٥. 1. ٥. 1. -1. 1. -1, 1. 0, -1, ٥. 2) (1, -2, -1, -2. 2. -3, 1, -2, 7) (1, 1, 0. -1, -1, ٥. 1. 1. 1)(1. 1. 0) C 1. 1. -2. 1. 1. 1. 1. 1, -1, -1, 3) (1, 1, 1, ٥. 1. ٥. 0)(1, 1. ٥, 7) (-1. -1. 1. ٥. 1. -1. 1. ٥. -1. 1. 1. 1. 0)(-2, -1, -1, 2. -2. ٥. -2. -1. 0, -1, 1. -1. -1. -1. 0. 3) (1, -1, -1, 0. 1. 0. 1. 1. 1)(0, -1, 1. 1. ٥, 1, -1, -1-2)(3, 1, 1, -2. -2. -2. 1. (1. -2, -2, 7) (1, 1, 0, -1, 1) (2, ٥, -2, -1, -1, -2, 4)(1, 0, -1, -1 -1. 3) (-2. -1. 2. -1. 0. -1. 1. ٥., -1. 0. 2. 1. 1. ٥. -1. 2. 1. 4 1, -2, 2, з, 0) (-1, -1, 1, 0. -1, -1, 0, -1, 4) (1, 1, 0. 1, 0, -1, -1, -1. 2)(1, -1, -1. 2) ٥. -1. 1. 1. 3)(1, -2, 3, -1. -1. 1. 1. ٥. 1. ٥. 1) (-1, -1, -1, -1, 1, 0, 0, 4) (1, 0, -1, -1, -1, 0. 1. -1. -2. 4. 4. 0) C 0. 1. з. ٥, 12)(0, 0, 0) (1, 4, -1, -2, -4, -4, -3, -2, 0. 0. 0. ٥, 0)(0,-1, 1, 30 (1. 1. 0. -1. 1. 1. 1. ٥. 1. 1. -1. ۰. -1. -1. ٥, 1. ٥. ٥. ٥. ٥. ٥. 0) (3, -1, з. -1. з. 2, -1, 2, 0)(1, 1, 1. 1. -2. 1. 1. 1. 0)(1. ο. 1. 1. -1. -1. 0. -1. 2) ٥, 4) (2, -2, 33 -1. 1. -1. -1. 1. -2. 1, -1, -1, -1. 5) (1, -1, 3. -1. -2. -3. -3. -2. 9)(1, ο. ٥. 1. -1. -1. -1. -1. -1. -1. 1. -1. 0, -1, 1. 3) (-1, ٥. 1. 2, 1, -1, 1. 2. 0)(1, 0, 1, -1, 0, -1, -1, 1. 2)(1, ٥. 2. -1, -2. -2. -2, -1, 6) ٥. 2. 1. 0) (0, -1, 0. ٥. 0) (1, -1, -1, 3)(0, -1, -2. 1. 2. 1. 1. 1. 1. ٥. 1. 0. -1. -1. ٥. 1. ٥. 1. 0) -1. ٥. -1. -1. ٥. 4) (-1, 0, 1. -1, ٥. -1. 1, -1, 3) (2, 1, 1, -1, -2, -1, -2, ٥. 4) (2, -1, 2. 2. 2, ٥. 1. -1. 2. 1. 0, -2, ٥. 0) (1, 0, ٥. ٥. ٥. ٥, ٥. 0)(1, ٥, ٥. 1) (2, -1, -3, 8) 1. 0. 1. -1. -1. -1. -3. -3. 1. 1. -2. -2. 1. 5) (-1, -1, -1 з, з, з, 0)(0, 1, 0)(1, -1. -2. 0. 0. 2. 1. 1. 0. -1, 0, 1. 1. 1. ٥. -1. ٥. 1. -1. 1. 0. -1. 1. 1) (-2, 1, 2. 4, 1, -2, 2, 3, 0) (1, -1, 4, -2, -3, -4, -4, -2, 12)(1, 1, ٥. -1. 0) ٥. 1. 1. 1. (1. -1, -1. 1, -1, 4) (-1, -3, 2, 1, -3, -2, -1, -3, 10) (3, 3, 0, -1, -1, -1, 0) (-1. -1. ٥. ٥. 2, 3, 0)(1, 1. ٥. ٥. -1. 1. 1. 1. 0. ٥. ٥. ٥. ٥. 1. ٥. ٥. 0) (0, -1, -1, 2, 2. 1. 2, ٥. 0)(4, 3, 1. -2, -2, 1. з, 0)(0, 2. -1. -1. -2. -2. 7) 4. -2. -1. 0 1. 1. 2, 1. -2. 1. 1. 0)(1, 1, 1, 1, 0, -1, ٥. ٥. 0) (0, -1, -1, 2, 2, ٥. 2, 1, 0)(-1, -1, 0, -1. ٥. 1. 31 -1. 1. -1. 2. ٥. 2, 2, -1, 0) (1, 3, -1, -1, -3, -3, -2, -2, 9) (1, -1, -1, 2, 1. 1, 1, -1, 1)(0, -1, -1, -1. 0, -1, 40 1. -1. 2. 1. -1, -1. ٥. 1. 2. 0) (1, -2, 2, -3, -1, -3, -3, 1. 9) (3, 4, 4, 1. -2, 0, -2, 1. 0)(0, 1. 1. ٥. 1. ٥. 1. -1. 0) 1. 1. 0. -2. 1. 1. 2. 0) (0, -2, 0, 1, 2, 1, 2, 1. 0) (1, -2, -1, -1. 2, -2, 0, -2. 6)(-2, 2. ٥. -2. -2. 1. -1. -1. 4) (-2, -1, -1, 7) (-1, 2, -3, -3, -1. -1. -1. 0. 1. -1. -1. 1. -2. 1, -2, -1, 1. -1, -2. -3. 10)(1. 2. -1. -1. 2. 2. ٥. 2. 0) (-3. -2, 2, -1, -2, -2, 1, -1, 8) (0, -1, 1, 0, -1, -1, -1, -1, 4) (1, -1, -1, 3, 2, 2, 2, -1. 0)(1,-1, 0, 1. 1. 1. 0. -1. 3, -2, -4, -4, -4, -1, 1, 11) (2, 0, 1, -2, -1, -2, -2, 1, 5) (-1, 1, -1, 0, 1, -1, -1, 0, (2. 3) (2, 3, 3, 2. 1. -4. 1. 1. 0) 1, 0, -1, 1, 0, -1, -1, 3)(-1, 1, 3, 1, -2, 1, 0)(0, 0, 0, 0, 0, -1, 1)(-2, 2, -1, 73 (-1. 1. 2. ۰. ٥. -1. 0. -2. -2. -1. 1. ٥. 0, -1, ٥. 1. 0)(1, 1. ٥. 0, -1, -1, -1, -1, 3) (1, 2, -2, -3, -3, -3, -1, 1. 9)(1,-2, 1, 1. 01 (1, 1. 1. 1. (-3. 1. 1. 1. 2. 2. 2, 1, 0)(1, 2. 2. ٥. -1, 1, -1, 1, 0)(1, -2, -3, -2. 1, -4, 3, -3, 10)(-1, 0, 1, 3) -1. -1. -1. C 0. -1. 1. 2. 2. 2, 1. -1. 0)(0, 1, 1. 1. ٥. 0, -1, 1, 0) (3, 2, 2, 2, -1, 0, -1, -1, 0)(-1, -1, -1, -1. ٥. 1. -1. ٥. 40 6 2. 2. 2, -1, з. -1. 3, -1, 0) (-1, 1, -1, ٥. ٥. -1, -1, -1, 4) (2, 1, 1. 1. -1. 0, -1, -1. 1) (2, 2, ٥. 1. -1. -1. -2, -2. 40 ٥. 0. 1. 1, 1, 1, -1, 0)(-1, 0, 0, ٥. ٥. 0. 0. 0. 1) (1, -1, 1. 1. 1. 2, -1, -1. 1)(0, 0. ٥. ٥. ٥. ٥. 1. ٥. -2. -2. -1. -1. 2. -1. ٥. 6) (1. 0. -1. ٥. -1. 1. 1, -1. 2) (0, -1, 1. -1. ٥, 1. 1. 1. 1)(0. -1. 1. -1. 1. 2. 0) 3, 3, -1, -1, -1, -1, з, 0) (-1, -1, -1, 1, -1, 0, -1, ٥. 4) (1, 1, -1, ٥. 1. 1. ٥. 0)(1, 1, 1, ٥. 0. 0. -1. 6 4. 1. 1. (-1, -1, 0, -1, 0, 1, 1. 1. 2) (1, 1, 0. -1, -1, -1, 0. 1. 2) (-2, 1. 2. 4. 2. -2. 1. з. 0)(-2, -2, 1. -1. -2. -1. 1. ٥. 6) ο, (0, -1, -1, -1, 0, -1, -1, -1, 0, 3, 3) (1, 1, 0) (-1, 0, 0)(1, 1. 1. 1. -1. 0. ٥, 1. 2. 2 1 -2. ٥. о. з, 1. 0)(-1, -1, 1. 0 -1. -1 1 ٥. 31 2. 0)(3, C 3. з. 2, -1, -1, -1, 1. 1. 1. -1. -1. -1, 2. 0) (2, -1, 2, ٥. 2. 2. -1. 1. 0) (-1, -1, 0, -1, 0, -1, 1, -1, 4)(-2, -1, 1, -1, -1, -1, 1, -1, 5)(1, 1. 0)(1,-1,-1, 0, 3) 1, -2, ۰. 1. 2. 2. -1. ۰. -1. 1. C 0. 1. 0. -1. -1. -1. -1. -1. 4) (2, 1, 1, 0, 1, 1, -2, 2, 0) (-2, 2. 1. 4. 2, -2, 1. з, 0) (1, -2, -2, -1, 1. -3. 2. -2. 7) 1, 1, -1, 1, 0) (1, -2, -1, -1, 2, 0, -2, -2, 6)(-1, ٥. ٥. 1)(2, 1, 0, (1. ٥. 1. 0. 1. 1. 1. -1. 1. -2. -2. -2. 51 C 1. -1. -2. -1. 0. -2. 2. -2. 6) (-1, 1, 1, -1, 0, ٥. 0. -1. 2)(1, 1. 1. 1. -1. ٥. ٥. ٥. 0)(1, -1, -1, -1. 1. 1. -1. 1. 1. 2. -1, -1, -1, 0, 0. 4) (3, 2, з, 0)(1, -1, 1. 1. -2. 0. 0) (-1, -1, 1. з. -1. 0. -1. -1. 1. ٥. 1. 0) (1, -2, -2, 2, 1, 0, 1, -1, 3)(1, -1, 0, 1, -1, -1, 1, -1, 3) (-1, 3, 3, -1, 2, 2, 2, -1, 0)(-1, 1, 0, 1. 1. -1. 0. 1. 1) (-1, -2, 1, -1, 1, -2, -2, 2. 6) (0, 1, -1, 1, -1, 1, 0, -1, 2) (0, 1, -1, -1, 2, 2, 1, 0)(-1, -1, 0, -1, 1, -1, -1, 2. 40 1. 10) (2, 2, -1, 0, (0, 0, 0, 0, 0, (2, -1, 2, 2, 2) ٥. 1. ٥. ٥. ٥. 0) (1, -3, -2, -2, 3, -4, 1. -3. -1. 2. 2. 1. 0)(1, -1, 2, -2, -1, -2, -2. ٥. 6) o) 2. 3, -1, -1, 0) (0, -2, -2, -2, -1, 2, -1, -1, 7) (3, -2, -3, 3, -1, -1, -3, 7)(-1, 1, 0, 2, 1, -1, 1. 1. 2. 0, 3)(1, -1, -1, 0, C 0. 1, -1, -1, 1. 1. 0, 1, 1)(1, 1, 0, 1, 0, 1, 1, -1, 0)(1, -1, 1, -1, 0, -1, -1, 2. 1. 2. 2. 01 0, 0, 0)(1, 0, -1, 1, 1, -1, -1, -1, 3)(-1, 1, 1, 2, 1, -1, 0, (0, 0, 1, 0, 0, ٥. 2. 0)

Method II

Convex hull computation

✓ Can we use less inequalities ? Yes!



Method II

Convex hull computation

Can we use less inequalities ? Yes!

Algorithm 1: Selecting n inequalities from the convex hull \mathcal{H} of an S-box

Input:

 \mathcal{H} : the set of all inequalities in the H-representation of the convex hull of an S-box;

 \mathcal{X} : the set of all possible differential patterns of an S-box;

n: a positive integer.

Output: \mathcal{O} : a set of *n* inequalities selected from \mathcal{H}

- 1 $l^* := None;$
- $2 \mathcal{X}^* := \mathcal{X};$

$$3 \mathcal{H}^* := \mathcal{H};$$

$$\mathbf{4} \ \mathcal{O} := \emptyset;$$

- 5 for $i \in \{0, ..., n-1\}$ do
- $l^* :=$ The inequality in \mathcal{H}^* which maximizes the number of removed 6 impossible differential patterns from \mathcal{X}^* ;
- $\mathcal{X}^* := \mathcal{X}^* \{\text{removed impossible differential patterns by } l^*\};$ 7

$$arepsilon$$
 $\mathcal{H}^* := \mathcal{H}^* - \{l^*\}$

 $\begin{array}{c|c} \mathbf{s} & \mathcal{H}^+ := \mathcal{H}^+ - \{l^*\};\\ \mathbf{g} & \mathcal{O} := \mathcal{O} \ \cup \ \{l^*\}; \end{array}$

10 end

11 return \mathcal{O}

Applications

- Automatic security evaluation with respect to single-key and related-key differential attacks.
 - obtain the lower bound of the number of active S-boxes of all characteristics
 - \checkmark useful in the design of block ciphers

Automatic search for single-key and related-key differential characteristics

- ✓ obtain characteristics with high probability
- useful in (related-key) differential attack, (related-key) boomerang attack, biclique attack ...

Application I : Security evaluation

- obtain the lower bound of the number of active S-boxes of all characteristics.
 - I. Set the objective function to be the sum of all variables indicating the activities of the S-boxes;
 - 2. Include the constraints imposed by the operations involved in the cipher;
 - 3. Require that all variables are 0-1;
 - 4. Solve the MILP model using the Gurobi optimizer , and the objective value of the optimized solution is a lower bound of the number of active S-boxes.



http://www.gurobi.com

Application I : Security evaluation

Iower bounds of the number of active S-boxes of the related-key characteristics of PRESENT-80

	With CDP Constraints		Without CDP Constraints	
Rounds		<pre># Time(in seconds)</pre>		
1	0	1	0	1
2	0	1	0	1
3	1	1	1	1
4	2	1	2	1
5	3	5	3	3
6	5	16	4	10
7	7	107	6	26
8	9	254	8	111
9	10	522	9	171
10	13	4158	12	1540
11	15	18124	13	8136
12	16	50017	15	18102
13	18	137160*	17	49537*
14	20	1316808*	18	<u>685372*</u>
15	<u> </u>	> 20 days	_	> 20 days

There is no related-key characteristic for 12+12=24-round PRESENT-80 with probability higher than $(2^{-2})^{16} \times (2^{-2})^{16} = 2^{-64}$

Warning!

Such bounds are only valid for characteristics, not for differentials

Application II : Characteristic search

obtain characteristics with high probability

- I. Set the objective function to be the sum of all variables indicating the activities of the S-boxes;
- 2. Include the constraints imposed by the operations involved in the cipher;
- 3. Require that all variables are 0-1;
- 4. Solve the MILP model using the Gurobi optimizer, extract a feasible solution when the objective value is small enough;
- 5. Check whether the solution is a valid characteristic. If it is invalid, add some valid cutting-off inequalities and go to step 4. If it is valid, we now have a characteristic.

Application II : Characteristic search

Improved 15-round single-key differential characteristic and differential for SIMON48, a lightweight block cipher designed by NSA.

SIMON48				
Rounds	Left	Right		
0	000000010000000000000000	00000100010001000000000		
1	0000000001000100000000	00000001000000000000000		
2	000000000010000000000	000000001000100000000		
3	000000000000100000000	000000000010000000000		
4	000000000000000000000000000000000000000	000000000000100000000		
5	000000000000100000000	000000000000000000000000000000000000000		
6	0000010000100000000000	000000000000100000000		
7	0000000001000100000010	00000100001000000000000		
8	001000010000100001000	000000001000100000010		
9	0000000001000100000010	00100000100001000001000		
10	0000010000100000000000	000000001000100000010		
11	000000000000100000000	00000100001000000000000		
12	000000000000000000000000000000000000000	000000000000100000000		
13	000000000000100000000	000000000000000000000000000000000000000		
14	000000000010000000000	000000000000100000000		
15	000000001000100000000	000000000010000000000		

The probability of the above characteristic is 2⁻⁴⁶; by considering the differential effect, the probability is 2^{-41.96}, which is the best result published so far for SIMON48.

Thanks!

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1.

2.

3.

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